Massively parallel computing on an organic molecular layer

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Modern computers operate at enormous speeds—capable of executing in excess of 10¹³ instructions per second—but their sequential approach to processing, by which logical operations are performed one after another, has remained unchanged since the 1950s. In contrast, although individual neurons of the human brain fire at around just 10³ times per second, the simultaneous collective action of millions of neurons enables them to complete certain tasks more efficiently than even the fastest supercomputer. Here we demonstrate an assembly of molecular switches that simultaneously interact to perform a variety of computational tasks including conventional digital logic, calculating Voronoi diagrams, and simulating natural phenomena such as heat diffusion and cancer growth. As well as representing a conceptual shift from serial-processing with static architectures, our parallel, dynamically reconfigurable approach could provide a means to solve otherwise intractable computational problems.

n current computers the logic can reconfigure itself for a new problem or even select a suitable circuit from a few available ones to evolve its hardware¹. However, once the logical path to solve a problem is determined, current is passed through a circuit that remains static. During serial computation, logical operations are performed in a strictly defined sequence to obtain the solution. In contrast, circuits of biological processors are dynamic during computation and all fundamental computing elements operate collectively and simultaneously. As an alternative to serial logic operation, von Neumann demonstrated parallel computing on a piece of graph paper by moving black and white dots together using simple rules²⁻⁴. To implement such a cellular automaton (CA) in hardware, each cell representing the dots should communicate with its neighbours simultaneously to generate a collective decision^{5,6}. Proposals are mounting for such network-based molecular computing7-9. Recently, suitable cells/dots have been built using quantum dots and molecules¹⁰⁻¹³. However, thus far, it has not been possible to assemble them into a two-dimensional grid where a CA cell communicates with many other devices at one time to execute a collective computation. In addition, in conventional cellular computing, the rules for updating cell states are fixed before computation, and the circuit or neighbourhood is kept static. Realization of dynamic CA circuits would enable us to address problems that are the prerogatives of natural bio-processors.

Here, instead of wiring single molecules/CA cells one-by-one, we directly build a molecular switch¹⁴⁻¹⁷ assembly¹⁸⁻²³ where \sim 300 molecules continuously exchange information among themselves to generate the solution⁵. We demonstrate the ability of a CA cell to change its neighbourhood from 2 to 6 CA cells in a controlled manner. Physically, it means that a molecule could interact locally with up to six molecules at a time during information processing. It should be noted that in the human brain, a neuron communicates with up to \sim 10⁴ neighbouring neurons at a time, and neural circuits evolve continuously during their life-span. By separating a monolayer from the metal ground with an additional monolayer, we developed a generalized approach to make the assembly sensitive to the encoded problem. The assembly adapts itself automatically for a new problem and redefines the CA rules in a unique way to generate the corresponding solution. We demonstrate the computing potential of the molecular assembly by realizing standard computational constructs such as logic gates, directed propagation of information, and so on. In addition, we have demonstrated the potential of dynamic circuit creation by physically encoding differential equations and generating quantitative solutions for heat diffusion^{24,25} and the mutation of normal cells into cancer cells²⁶. The molecular assembly functions similarly to the graph paper of von Neumann, where excess electrons move like coloured dots on the surface, driven by the variation of free energy that leads to emergent computing^{27–30}. The assembly processes ~300 dots at a time, whereas the fastest processors today operate only on one bit/dot at a time per channel.

Realization of a DDQ computing grid

The 2,3-dichloro-5,6-dicyano-*p*-benzoquinone (DDQ) molecules (Fig. 1a) that reversibly switch among the four conducting states^{17,18} 0, 1, 2 and 3 are deposited on an Au(111) surface as a bi-layer (Supplementary Sa). To encode information, we select molecules one-by-one by physically moving the scanning tunnelling microscope (STM) tip to the highest current location on the molecule, as shown in Fig. 1b. Then the desired conducting/logic states 0, 1, 2 or 3 are written by applying a pulse of -1.6, 1, 1.3 or 1.6 V tip bias respectively for $5-10 \,\mu$ s at $\sim 300 \,\text{K}$. State 3 appears as a brighter sphere in the STM image due to having one extra trapped electron than in state 1. State 2's brightness is between that of state 1s and state 3s as it has a more elongated prolate-shape than state 0. Figure 1c shows the detection of the four states 0, 1, 2 and 3, which are denoted by blue, green, yellow and red balls respectively (Fig. 1d).

The connecting region between two neighbouring DDQs with more than 60% of the peak tunnelling current observed on the molecules is considered as a wire, which is represented by a solid line in Fig. 1e and f. Thus, we get a circuit where one molecule is connected to a distinct number of neighbouring

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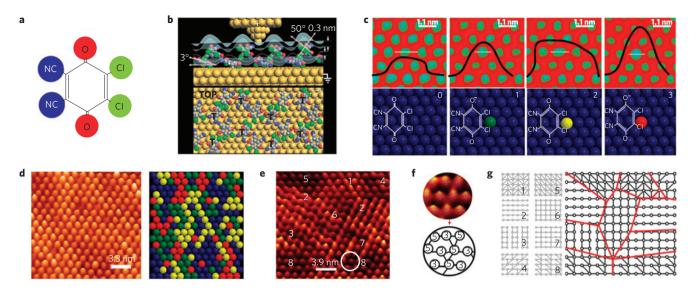


Figure 1 | **The concept of a wireless molecular circuit. a**, The DDQ molecule. **b**, The DDQ bilayer's atomic structure; side view (above); top view (below, Supplementary Movie S1). T denotes a molecule on the top layer. **c**, STM images for the four states (top); corresponding ball models and atomic structure (bottom). Current profiles of the four states (black line) taken along the white lines having current heights of (at 0.68 V) \sim 0.15 nA (state 0), 0.25 nA (state 1), 0.28 nA (state 2) and 0.32 nA (state 3) respectively. The height profiles are \sim 0.9 nm wide. **d**, A 13 × 17 matrix (\sim 400 nm²) and its ball representation (right). **e**, STM image at 0.2 V tip bias and 0.05 nA tip current. **f**, Zoomed region of the STM image and the corresponding circuit (below). The black rings represent molecules; the number represents interconnecting lines (wires). **g**, Eight possible circuits (left) in image **e** are detected (right). The red lines are Voronoi cell boundaries.

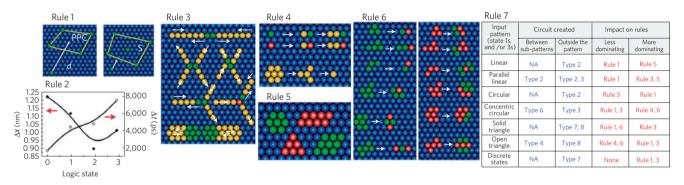


Figure 2 | **Discrete logic-state transport rules.** Rule 1: Convergent universe. A charge moves a distance *d* towards PPC (left). To calculate PPC, we neglect those areas that have charge(s) in between (right). Rule 2: Creation of spatial Δx and temporal limit Δt : the $\langle \Delta x \rangle$ (experiment) and $\langle \Delta t \rangle$ (simulation) are plotted for four surfaces initially covered with state 0s, 1s, 2s, or 3s. Rule 3: Divergent Universe. Examples of collisions; arrows denote the direction of motion of logic-states. Rule 4: Life of logic states. Trail of state 2 for motion of state 1 and 3 (top); the death of a cluster of state 2 (bottom). Rule 5: Collapse of space. Four examples of group formation of state 1s, 3s. Rule 6: Transformation. Fusion of two state 1s to create a state 3 (left); breaking of state 3 to create two state 1s (right). Rule 7: Priority of rules. Correlating input patterns, initial circuits and the dominant rules (see Supplementary Movie 4 for details and Supplementary Information text online for the algorithm to program these rules).

molecules with which it interacts at a time. In this article this interaction is referred to as one-to-many interaction at a time. The neighbourhood of a molecule or CA cell varies from 2 to 6 in the eight molecular circuits possible in the top monolayer. Continuous scanning of the bi-layer, by changing the tip bias from -2V to +2V in a loop, reveals eight distinct circuits (Fig. 1g). In a matrix, if states 0 or 2 are in excess of 60% within $\sim 20 \text{ nm}^2$ area, the DDQs prefer to reassemble into a circuit of type 7 or 2 respectively. If the number of state 1s are more than 50% or the number of state 3s are more than 30% in a $\sim 20 \text{ nm}^2$ area, then a circuit of type 1 or 5, respectively, is created (Fig. 1f,g). As we write a matrix, depending on the concentration of excess electrons (state 1s and 3s), DDQs in some part of the matrix region re-orient to the nearest energetically favourable circuit. Hence, for every new input matrix, a unique arrangement of several circuit-domains is created automatically

(Fig. 1g, Supplementary Movie S3). Thus the density of free electrons in an input pattern, the transformed circuit and the logic state transport rules are correlated.

Switching of an entire molecular arrangement as a function of applied bias has been reported for different molecules (Supplementary Sb; refs 19–23). However, the coexistence of multiple circuits side-by-side, as demonstrated in Fig. 1g, has not been observed. Owing to weak inter-layer coupling (Fig. 1b) and strong coupling among surface molecules, the top monolayer relaxes almost independently. This leads to the survival of several circuits. Therefore, excess electrons supplied to the assembly through states 1 and/or 3 find themselves in a potential surface of valleys and hills⁵; this triggers their spontaneous motion to minimize the free energy. As the encoding process re-defines the mode of interaction of DDQs with their neighbours, the excess electrons do not move randomly, but rather they follow well-defined rules (Supplementary Sc). In

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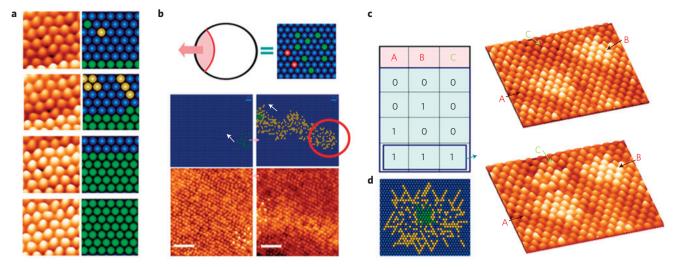


Figure 3 | **Information encoding, retrieving, transport and logic-gate operation. a**, STM image (left) and its equivalent ball representation (right) are shown in four steps (top to bottom) of an event where all state 0s in a 6×7 grid are converted into state 1s. **b**, Top row—schematic representation of an information packet. The shaded region has a higher electron density, which determines the direction of motion. A real packet is shown to its right. Middle row—a simulated transport of an information packet comprised of a few state 1s (left-input). Yellow balls trace the transport path for 2,000 steps (right-output). Bottom row—STM images for initial (left) and final steps (right) of information transport (scale bar: 7.5 nm). **c**, Truth table for an AND gate (left). STM images (right; scale bar: 2.8 nm) show A = 1, B = 1, C = 0 (top-input) and A = 1, B = 1, C = 1 (bottom-output). **d**, AND gate operation with (A = 1, B = 0) and (A = 0, B = 1) should not have evolved state at location C; A or B should retain logic state 1. 35 state 1s used to create A; simulation shows that state 1s collapse to maintain the logic state 1 (A = 1).

particular cases, the spontaneous evolution of the input matrix requires an extra trigger by scanning the surface at -0.98 V. All states of the entire logic pattern are erased to state 0 by scanning the surface at -1.6 V, and it is possible to use the same surface repeatedly for pattern evolution.

Logic-state transport rules

We have identified the following seven categories of rules (Fig. 2) that dictate logic-state transport.

Rule 1: if negatively charged state 1s and 3s are distributed heterogeneously in an input pattern, a pseudo-positive charge (PPC) is generated to which they are attracted from at most ~15 molecules apart. As a result, state 1s and 3s move to this single point. As ~700 molecules influence each other at a time, it is not feasible to express this rule in an abstract form; it would require ~4⁷⁰⁰ rules. Therefore, we adopt an analytic approach to program this rule (Supplementary Sh). The dynamic feature of the rule is provided in Supplementary Movie S4.

Rule 2: a state 2 site offers a lower barrier for the electron transport than a state 1 site, and a state 1 moves faster than a state 3. By tuning the distance Δx between two DDQ molecules (0.93–1.03 nm) in an assembly of circuits we can speed up or slow down the electron transport process as described in Fig. 2 (Supplementary Sd). To program this rule in our simulator such that the temporal behaviour of transition-state dynamics is consistent with our experiment, we leave a fraction of states 1 and 3 unchanged in each update.

Rule 3: state 1*s* and/or 3*s* and their groups experience repulsive forces in a particular direction if they come in contact at a single point. This rule promotes divergence of bits in a logic pattern.

Rule 4: state 1*s* and/or 3*s* and their groups move leaving a trail of state 2*s* that remain static. If seven sets of state 2*s* form a hexagonal cluster, then all switch to the lower energy state 0 in two steps, except the central DDQ.

Rule 5: in particular arrangements, clusters of state 1s and/or 3s form a group and move on the surface as a single entity. Groups lose their property of a single entity if the contact dimension is at least two molecules in length.

Rule 6: when negatively charged states form a cluster that is asymmetric, the sets of cluster changes to a symmetric shape through the creation of two sets of state 1*s* by breaking one state 3 or through the creation of a new state 3 by fusing two state 1*s*.

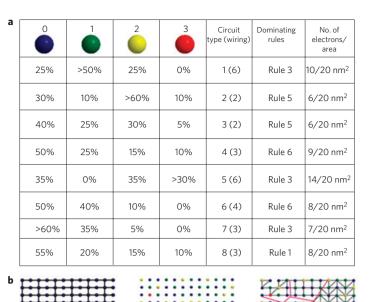
Rule 7: depending on the charge density of an input pattern, a unique composition of circuits is formed. In a CA grid, a particular set of rules is favoured in a typical circuit-domain (see tables in Figs 2 and 4). The order of relative circuit areas for an input pattern determines the order of execution of the rules. The algorithms used to program these rules in our simulator are described in the online text (Supplementary Sh).

Conventional CA computing

Writing, erasing and retrieving information. In Fig. 3a we demonstrate the sequential writing of a state 1 matrix on a state 0 surface. The states are stored as static information until spontaneous pattern evolution is triggered externally. By scanning the surface at -1.68 V one can reset all molecules to state 0, thus erasing the information. To retrieve information the surface is scanned at ~ 0.2 V (Fig. 1d).

Directed propagation of information. To send a complex information packet in a particular direction, we need to write additional states 1*s* and/or 3*s* so that an electron density gradient is created along the desired direction of propagation (Fig. 3b top). One example of transport over an apparently unbounded distance is shown by means of simulation and experiment in Fig. 3b (bottom). However, it is possible to send a packet to a particular location. The signal propagation stops when the motion of states within the packet generates a pattern that periodically repeats in a fixed space. To send an information packet that is in the form of a group, we need to create its mirror image in contact with it; repulsion would then trigger the propagation of information following rule 3, which dominates in the circuits 1, 5 and 7 (Fig. 4a).

Logic gate. The effects resulting from rule 3 appear similar to the interactions in the billiard ball model³¹, which has been used to design logic gates. We have realized an AND logic gate based on



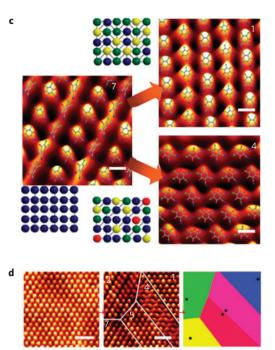


Figure 4 | Computing constructs, density classification and Voronoi decomposition. a, Relations among density distribution, circuit type and dominating CA rules. b, Schematic representation: all CA cells in state 0 (circuit 7, left); the encoded input pattern (middle); the derived composition of circuits (right). c, STM image of a CA grid (6 × 5) comprised of state 0s (circuit 7), ball representation below. Two patterns (coloured balls, right) are encoded. One pattern converts the grid to circuit 1 (top) another to circuit 4 (bottom). Scale bar is 1 nm. d, 13 × 15 CA grid at state 0 (left); a pattern is encoded that converts it into a composition of circuits (middle) noted by circuit number. Voronoi decomposed cells containing the Voronoi points (black dots) derived from the middle image (right).

interactions in which rule 1 dominates over rule 3. A schematic logic device is shown in Fig. 3c. A random composition of states 1 and 3 (density > 0.5 electrons nm⁻²) written in a circular form is a logical '1' and the absence of any such composition is a logical '0'. If we write two logical '1's at most 15 nm (15 cells) apart, only then is a new composition created automatically on the surface depicting logical state '1'. If any one or both input compositions are absent we do not get such an output; rather states collapse at the same space (Fig. 3d), we get a logical '0' at the output location. Thus we realize an AND gate. A large number of such logic gates could be operated in parallel on DDQ CA by separating those by ~15 CA cells. The output of a logic gate could be transported to the input of another logic gate, as described in Fig. 3b.

Density classification. Density classification (DC) is a task to categorize particular elements in a mixture. DC is challenging when elements of different densities interact and influence each other at the same time. Classically, in CA, minority states in a space are converted into the majority's state to reflect classification³². In our CA, however, we divide local space into several strictly defined regions on the basis of charge density, as shown in the table of Fig. 4a. We can categorize the region into eight circuit classes with a maximum sensitivity of 1 electron/20 nm² (Fig. 4b). We have converted a state 0 space into two distinct circuit patterns simply by writing two input patterns with distinct charge densities. The CA grid spontaneously classifies two distinct density distributions with a difference of ~2 electrons/30 nm², by converting the grid into two distinct circuit domains (Fig. 4c).

Voronoi decomposition. In a Voronoi decomposition³³ a space is divided into multiple regions; each of which corresponds to a point such that any two neighbouring points are equally separated from the boundary between them. As described earlier, an organic monolayer when imaged at a low bias (~ 0.2 V) reveals the composition of circuits generated by the input pattern. The borders between the circuit domains are remarkably linear (see Figs 1e and 4d). The density classification at larger scales (\sim 50 nm²) triggers segmentation of the computing space and decomposes the surface into Voronoi cells. A physical process in the molecular assembly could be modelled in terms of the evolution of the points (black dots in Fig. 4d) representing the Voronoi cells.

Mimicking two distinct natural phenomena

Biological computers like the brain do not have logic gates yet they solve complex problems. During computation, the encoded information pattern in the neurons dynamically modulates the neural architecture and continuously evolves to reach a collective solution. Theoretically, by tuning input patterns and CA rules, we can solve several problems without using any logic gates³⁴. Here we mimic two natural events in the molecular CA matrix by tuning input patterns and effective CA rules. We encode two distinct input patterns that evolve over time in such a way that the transport of free electrons (or logic states) follows the essential features of diffusion^{24,25} for one input pattern and follows the evolution of cancer cells²⁶ for the other input pattern. We have also emulated the pattern evolution in a simulator. The global features of the experimental patterns are in reasonably good agreement with the simulated patterns (Supplementary Se).

In the diffusion process, a blue ball (state 0) denotes a normal material. When it accepts one electron, it turns to green (state 1). Two excess electrons turn blue to red (state 3). When electrons leave, green and red balls relax to state 2s, which are yellow (Fig. 5a). To create a directional flow, we write a pattern of straight-lines each composed of alternating state 3 and state 1 so that they form a group and move as a single unit (Fig. 5b). The direction of flow is controlled by the potential gradient of states in the background^{24,25}. For free-energy minimization, the linear arrangement tries to bend into a circular shape, but as it

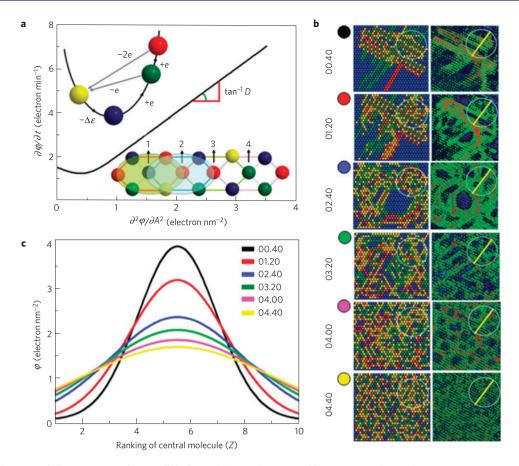


Figure 5 | **Mimicking natural phenomenon I: electron diffusion. a**, Electron flux rate $\partial \varphi(t)/\partial t$ versus surface-electron concentration variation $\partial^2 \varphi(t)/\partial A^2$. Inset (top)—transition scheme of DDQ logic-states. Inset (bottom)—a linear array of unit cells and their ranking (1, 2, 3, 4). The shaded green and blue regions are unit cells. **b**, STM images (24 × 27 nm) at 0.68 V tip bias and 0.05 nA tip current (right) with the corresponding simulation (left) shown side-by-side with time (minute) (Supplementary Movie S5). **c**, Temporal (minute) variation of electron concentration φ along a linear array of 10 unit cells (*Z*: 1–10) in the white ring of **b** is plotted.

begins to move, the coupling breaks. The broken parts follow the potential gradient independently (Supplementary Movie S5). We calculate the concentration variation and the rate of electron flow as described in the Methods section and plot them in Fig. 5a. A linear relation suggests a diffusion process; its slope D (2 nm² min⁻¹) is the electron diffusion coefficient. Figure 5b shows that gradually over time states 1 and 3 spread homogeneously on the surface. From Fig. 5c we get flux $\varphi(x, t) = 14/t^{1/2} \exp(-(x-5.5)^2/t)$ as the quantitative formal solution of the diffusion equation on a linear array of ten cells.

In a conventional CA, as transition rules are fixed, irrespective of the input pattern one can generate only one kind of dynamics of logic states^{3–5}. To demonstrate the potential of CA rule modulation during computation (Rule 7, Fig. 2), we generate a significantly different type of dynamics in the same CA grid by emulating the mutation of normal cells to cancer cells. To formulate the mutation process, we correlate four transition states of the normal cell to four conducting states of a DDQ. A normal cell (state 0) first mutates inactivating a single tumour suppressor gene (TSG) (state 1) at a rate u_1 ; it then mutates again inactivating another TSG (state 3) at a different rate, u_2 (Supplementary Movie S5, Fig. 6a). Cellular proliferation occurs when one state 3 creates two cells in state 1. Depending on the input pattern, the fusion of two state 1s into one state 3 may also dominate. When cancer spreads in tissues, it may leave a trace of chromosomal instability (CIN) or state 2. In theory, the inactivation of TSGs follows three successive differential equations²⁶. The solution for the third equation (see Methods section) provides the number of cancer cells (state 3s) produced,

which in our CA grid is the count of new red balls generated (N3). N3 depends on the cell population, N, in a tissue compartment (Fig. 6b). We write state 1s in the form of two concentric rings so that they neither converge nor diverge (Fig. 6c). Thus, an artificial tissue boundary CG is established, wherein N is the total number of DDQs inside CG. We tune u_1 by switching some of the state 0s to state 1s inside CG between every two scans, whereas u_2 is decided by the system itself.

We change *N* by modifying the separation between the two rings while keeping the inner ring unchanged. Here, the half-life $t_{1/2}$ for state 3 is the duration for which *N*3 is increased to 2*N*3. We plot average $t_{1/2}$ values for different values of *N* in Fig. 6a. A similar feature between the half-life $t_{1/2}$ versus *N* plot in Fig. 6a and the kinetics of cancer²⁶ suggests a consistent encoding of a two-rate controlled phenomenon in the CA grid.

Figure 6c shows that, at a very low population ($N < 290 \sim 1/\sqrt{u_2}$), when the two rings are less than 3 nm apart, the inner ring stabilizes by rearranging the states, after which both the rings merge. In this case, almost every collision between two state 1s produces a state 3, and N3 varies with time as $\sim 4t^2$ (Fig. 6b). The conversion continues at this rate until all state 1s are converted to state 3s. For an intermediate population (290 < N < 620), the intercluster distance increases to ~ 5 nm, and the state 1s of two rings collide before reaching an equilibrium. Instantly produced state 3s break into state 1s, and N3 varies as $\sim 18t$. For a very high population ($N > 620 \sim 1/u_1$), when two rings are more than 8 nm apart, a higher PPC leads to a large-scale one-to-one collision of state 1s; thus conversion to state 3 is accelerated. Therefore, the

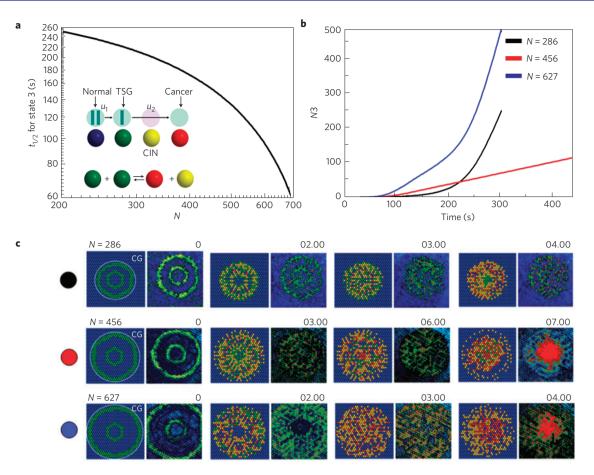


Figure 6 | Mimicking natural phenomenon II: evolution of cancer cells. a, Measured half-life $t_{1/2}$ (seconds) to produce state 3 versus tissue population *N* on a log-log scale. First row (inset)—a real cell with green lines as TSG with corresponding ball representation (second row). Third row—a scheme for 2 one-TSG-deactivated cells creating one cancer cell and one CIN cell. **b**, Number of state 3s produced (*N*3) versus time for different *N*. **c**, Simulated (left) and corresponding STM images (24×27 nm) at 0.68 V tip bias and 0.05 nA tip current (right) are paired and ordered in a row with time (min). Simulations with N = 286, 456 or 627 having 149 (~50%, top), 196 (~40%, middle), or 222 (~34%, bottom) state 1s respectively. After each STM scan of 40 s, 5 (top), 8 (middle), and 11 (bottom) new state 1s are written inside the respective CGs within 20 s (Supplementary Movie S5).

abundance of cells in state 1 allows N3 to vary as $\sim 40t^2$. N3s in Fig. 6b provide quantitative solutions for the differential equations. During evolution, if state 2s are erased or written, the production of state 3s is also slowed down or accelerated respectively; this is a distinct feature of CIN dynamics. As N3 kinetics changes with time, different sections of the lines in Fig. 6b follow additional solution functions $7e - 4t^2$, $7e - 10t^3 \dots t^5$ for low populations and $0.05t^2$, t^2 , $t^{1.8} \dots t^{1.2}$ for high populations. Remarkably, these smaller parts reveal complex CIN dynamics, which are consistent with the established models²⁶.

During realization of diffusion and cancer growth on a 24 × 27 molecular matrix, changes in the STM image contrast at 300 DDQ sites reveal that the rules are executed at least at 300 sites simultaneously¹⁸. Cancer may evolve naturally in real tissue at $\sim 10^9$ cells in ~ 100 yr, and heat diffuses in metals at a speed of ~ 100 km s⁻¹. However, in a molecular assembly, ~ 600 cells mutate in 6 min, and heat diffuses ~ 10 nm in 5 min. By writing an input matrix, we can tune the grids Δx and hence Δt , and solve several differential equations similar to conventional methods.

Outlook

To realize a CA that can carry out a wide variety of computational tasks, it is necessary to obtain a sufficient level of control on the transition state dynamics. The DDQ CA in this article provides an intriguing avenue to achieve such control while still relying on a relatively small and simple molecule. It is not only the interactions of the molecules in the DDQ CA, but also the subtle formation of circuits in the molecular top-layer that count in facilitating transitions between states. The evolution of circuits is dependent on an easy-to-control parameter: the charge density in an area. These circuits dramatically influence the dominance of transition rules, as we have observed, and offer an efficient way to influence the computational behaviour of the CA. The robust functioning of local circuits originates from the CA cell's one-tomany communication and interaction at a time^{18,30}. Generalization of this principle would change the existing concept of static circuitbased electronics and open up a new vista of emergent computing using an assembly of molecules^{6,27–29}.

Methods

Correlating diffusion parameters with the assembly states. The rate of change of electron flux φ is proportional to the variation of the gradient of electron flux, or $\partial \varphi / \partial t = D \nabla^2 \varphi$; *D* is the diffusion coefficient. Here, $\nabla^2 \varphi \sim \partial^2 \varphi / \partial A^2$ where $A (\sim 2.4 \text{ nm}^2)$ is the area enclosed by seven DDQ molecules (unit cell, Fig. 5a).

In a unit cell, φ is the ratio of total excess electrons Q(= number of state $1+2 \times$ number of state 3) and A. Now, $\partial \varphi / \partial A \approx \Delta \varphi / UC$; $\Delta \varphi$ is ΔQ between two adjacent unit cells, say 1 and 2 (or 3 and 4) of Fig. 5a (inset, bottom), and UC is the non-overlapped region of two unit cells (~1.3 A). The difference in $\partial \varphi / \partial A$ between two neighbouring unit cells, say 1 and 3 of Fig. 5a (inset), is calculated as $\partial^2 \varphi / \partial A^2$ and the value is assigned to the centre of unit cell 2 where $\partial \varphi / \partial t$ is measured. Here, $\partial \varphi / \partial t \approx \Delta \varphi / \Delta t$; Δt (40 s) is the time lapse between two STM scans.

Similarly, for every unit cell of the 24 × 27 matrix, we calculate $\partial^2 \varphi / \partial A^2$ from an STM image at $t = t_0$ and $\partial \varphi / \partial t$ from two STM images at $t = t_0$ and $t = t_0 + \Delta t$.

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For each value of $\partial^2 \varphi / \partial A^2$, the corresponding values of $\partial \varphi / \partial t$ are plotted in Fig. 5a. We have also plotted φ versus ranking Z (0, 1, 2, ..., n) of the central molecule of the unit cells along a straight line in Fig. 5c. Therein, the electron flux diffusion saturates in 4 min 40 s (noted as 4.40). We have taken the weighted average of flux to convert the staircase characteristic into a continuous trace in Fig. 5c; this is essential to find a formal quantitative solution. Equations of higher order and/or degree can be solved similarly.

Correlating cancer-cell evolution parameters with the assembly features. We have three coupled differential equations and two variables: states 1 and 3. As state 1s cannot be generated spontaneously, we encode state 1s using STM between two scans at a rate of $\sim 1.75 \times 10^{-2}$ min⁻¹ so that a mutation rate of u_1 ($\sim 1.61 \times 10^{-3}$ min⁻¹) for state 0 is maintained. The evolution of state 0 follows $dX_0/dt = -u_1X_0$, where $X_0(t)$ is the number of effective state 0s at time t, given by $X_0(t) = X_0(0) \exp(-u_1 t)$. Here, $X_0(t=0) \approx$ (initial state 1s in the rings + externally encoded state 1s), which is the number of effective normal cells (state 0s) involved in generating state 3s. The state 1s that leave the rings and the state 1s encoded by means of STM to compensate for the loss during state 3 production contribute together to the evolution of state 1; this is described by $dX_1/dt = u_1X_0 - Nu_2X_1$; its solution is $X_1(t) = X_0(0)[u_1(e^{-u_1t} - e^{-Nu_2t})/(Nu_2 - u_1)]$. State 1 mutates at a rate u_2 to produce X_3 number of state 3s following $dX_3/dt = Nu_2X_1$, which yields $N3 = X_3(t) = X_0(0)[1 - (Nu_2e^{-u_1t} - u_1e^{-Nu_2t})/(Nu_2 - u_1)]$. If state 2s are deleted between two scans, u_2 decreases to $\sim 1.18 \times 10^{-5} \text{ min}^{-1}$ from $\sim 3.96 \times 10^{-2} \text{ min}^{-1}$ for N = 627. Thus, CIN increases u_2 by $\sim 10^3$ times, which prompts us to simplify $X_3(t)$. For a smaller timescale that is considered in our experiment, a small population yields $N3 \sim X_0(0)Nu_1u_2t^2/2$; an intermediate population yields $N3 \sim X_0(0)Nu_1\sqrt{u_2}t$; and a large population yields $N3 \sim X_0(0)Nu_1u_2t^2/2$; here $u_1 \sim 1.61 \times 10^{-3} \text{ min}^{-1}$, $u_2 \sim 1.18 \times 10^{-5} \text{ min}^{-1}$ (ref. 26). The number of new red balls, N3, is obtained by comparing two consecutive STM images. The average values of N3 obtained by repeatedly evolving the same input patterns, as demonstrated in Fig. 6c, are plotted in Fig. 6a,b.

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Author contributions

A.B. designed the research; A.B. did the experiment; A.B. developed the CA simulator; R.P., A.B. and S.S. did the theoretical studies; A.B., R.P., S.S. and F.P. analysed the data; and A.B., R.P., F.P. and S.S. wrote the paper together; D.F. reviewed the work.

Additional information

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